

**ECE 436**

ADVANCED DIGITAL DESIGN

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# **Jackal Version 3.0**

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**A 16-BIT REDUCED INSTRUCTION SET COMPUTER SPECIFICATION FOR THE ECE436  
ADVANCED DIGITAL DESIGN COURSE AT THE UNIVERSITY OF VIRGINIA**

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# Jackal Version 3.0

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## INSTRUCTION SET INTRODUCTION

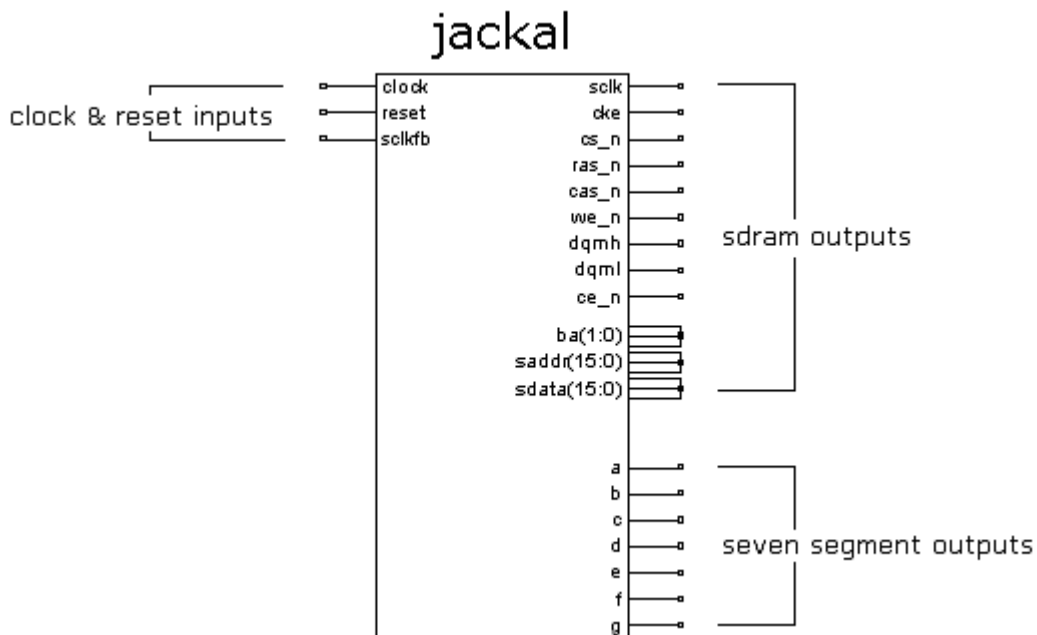
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### OVERVIEW:

The Jackal is an academic general-purpose reduced instruction set computer (RISC) architecture capable of hosting a wide variety of computational applications. This particular instruction set architecture (ISA) holds to the strictest definition of RISC by providing a very thin core of 16 possible instructions. The ISA's breadth ensures that a student team can design, simulate, synthesize, and place-and-route the processor core to an FPGA prototyping platform in the period of one semester.

A novel aspect of the Jackal ISA lies in its flexibility. While the 16-bit instruction set allocates up to 16 instructions and 16 possible registers, the specification identifies the functionality of 12 of each. This incomplete specification allows students to research and develop a strategy for defining the function and interface to the remaining four instructions and four register operands.

This processor is an educational tool for demonstrating the hardware/software interface while encouraging students to apply skills learned from previous computer engineering courses. As an educational tool that emphasizes simplicity, however, some desirable general-purpose processor (GPP) features have been sacrificed. First, this processor only has a 128KB address space and has limited I/O capabilities. A second deficiency with this processor's architecture is in its lack of interrupt handling capability. Finally, this specification does not include a floating-point unit (FPU); however, this absence does not limit this processor's ability to execute a variety of mathematically-involved algorithms. All students will have to keep cognizant of these deficiencies when writing applications for this processor core. Below, a high-level schematic capture of a previously implemented core depicts the inputs and outputs of the Jackal architecture when mated to an XESS XSA-100 FPGA prototyping platform (see Appendix).



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## INSTRUCTION SET OVERVIEW

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### SEMANTICS:

The bit numbering scheme will always require that bit strings read from right to left starting with bit 0. The core is little-endian, meaning that bytes at the lowest addresses have the lowest significance. The instruction word size is two bytes or 16-bits. Integer data is represented by two's complement form.

### INSTRUCTIONS:

The Jackal uses a 4-bit opcode (allowing for a total of  $2^4$  instructions) to reference each of the instructions; however, only 12 instructions have been defined. The remaining four instructions may be defined by the student designers. The 12 instructions have been selected to ensure that the core is fully capable of general-purpose processing. Additionally, the core's 12 instructions represent a cross-section of the most frequently used instructions for a general-purpose RISC processor.

### REGISTERS:

The Jackal is capable of referencing a total of  $2^4$  registers with its 4-bit register referencing operands. To permit a greater design space, however, only 12 of the registers operands have been formally assigned to general-purpose registers (R0 to R11). The role of the remaining four undefined operands (U0 to U3) is a task delegated to the student designers.

OPERAND	ROLE	OPERAND	ROLE	OPERAND	ROLE	OPERAND	ROLE
0000	R0	0100	R4	1000	R8	1100	U0
0001	R1	0101	R5	1001	R9	1101	U1
0010	R2	0110	R6	1010	R10	1110	U2
0011	R3	0111	R7	1011	R11	1111	U3

Two addressing modes are available on the Jackal core, register-indirect and immediate. With register-indirect addressing, a register is loaded by using the 16-bits stored in another register to provide the address of data stored at that memory location. With immediate addressing, the low and high bits of a register are loaded through the fixed value in the immediate instruction.

The program counter is a 16-bit register that may be addressable through the 4-bit operand-addressing scheme if the student designers so choose. Resets are triggered logic low and set the data in the PC and registers to 0x0000. Three single-bit condition registers, CRN, CRZ, CRP, store the output of the CMP instruction.

### MEMORY:

The Jackal's 16-bit memory address space, is capable of referencing  $2^{16}$  memory locations. Each memory location consists of 2 bytes (16 bits). The memory addresses range from location 0 (0x0000) to 65,535 (0xFFFF). All reads and writes to and from memory are accomplished through the LD/ST instruction. Programs written for the core will start at memory address 0 (0x0000) and may not exceed the range of the address space.

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## INSTRUCTION SET PREVIEW

### INSTRUCTION SET CONVENTION:

To minimize the confusion of describing the operation of each instruction, a convention has been established to help with the understanding of this ISA.

CONVENTION	EXAMPLE	DEFINITION
0x	0xFFFF	Denotes Hexadecimal Number FFFF
REG[ ]	REG[SOURCE]	Data in Register Referenced by SOURCE
MEM[ ]	MEM[SOURCE]	Data at Memory Address in SOURCE Register
VAL[ ]	VAL[OFFSET]	Value Denoted by 8-bit OFFSET

### INSTRUCTION SET SUMMARY:

The table below summarizes the instructions explicitly defined by this ISA. To promote simplicity in the ISA, several of the instructions perform a dual purpose. For example, register-to-register data movement can be accomplished with the AND or OR instruction when the SOURCE TWO register contains the data 0xFFFF or 0x0000 respectively. The NAND instruction can act as a logical negation (NOT) when the data at SOURCE TWO is 0xFFFF. The branch functions as a PC-relative jump instruction if the correct MODE applies. Thus, to write efficient assembly routines for this core, student teams will need to explore the nature of each instruction in detail. Further information about each of the instructions can be gained from the latter portion of this document. Please note the presence of undefined instructions (UDI) and undefined operands (UDO). These capabilities may be provided by the students.

INSTRUCTION	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ADD	0	0	0	0	DESTINATION				SOURCE ONE				SOURCE TWO			
SUB	0	0	0	1	DESTINATION				SOURCE ONE				SOURCE TWO			
AND	0	0	1	0	DESTINATION				SOURCE ONE				SOURCE TWO			
OR	0	0	1	1	DESTINATION				SOURCE ONE				SOURCE TWO			
NAND	0	1	0	0	DESTINATION				SOURCE ONE				SOURCE TWO			
SLA	0	1	0	1	DESTINATION				SOURCE				OFFSET			
SRA	0	1	1	0	DESTINATION				SOURCE				OFFSET			
LD	0	1	1	1	DESTINATION				SOURCE*				MODE			
ST	0	1	1	1	SOURCE				DESTINATION*				MODE			
LIL	1	0	0	0	DESTINATION				IMMEDIATE							
LIH	1	0	0	1	DESTINATION				IMMEDIATE							
CMP	1	0	1	0	UDO				SOURCE ONE				SOURCE TWO			
BRN/BRZ/BRP/JMP	1	0	1	1	MODE				OFFSET							
UDI	1	1	0	0	UDO											
UDI	1	1	0	1	UDO											
UDI	1	1	1	0	UDO											
UDI	1	1	1	1	UDO											

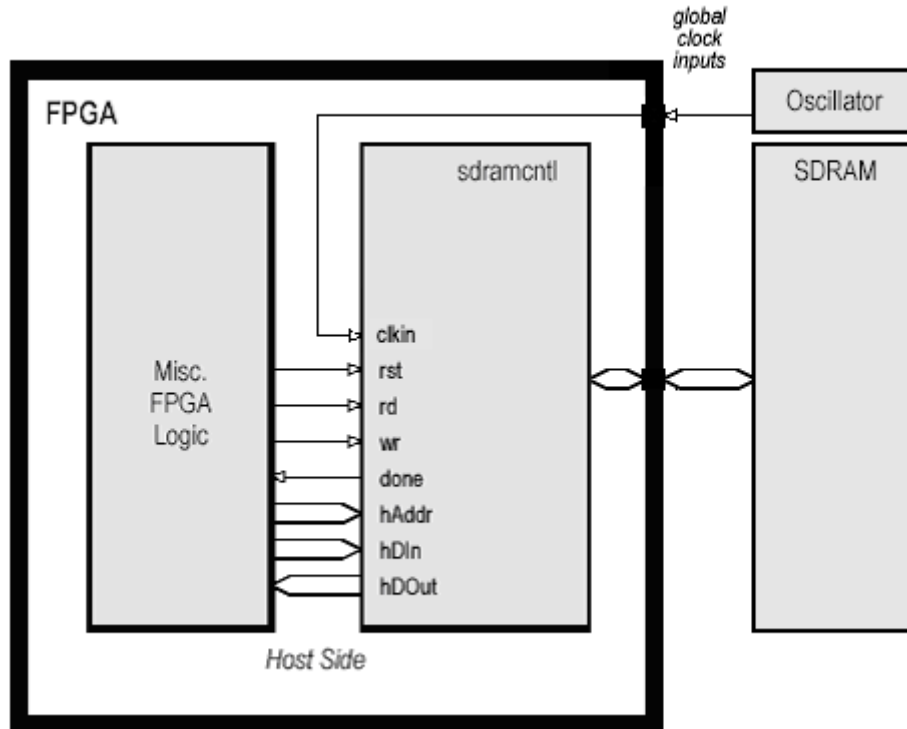
\*For the LD/ST instruction, the source/destination operand is a register holding a memory address.

# Jackal Version 3.0

## INSTRUCTION SET MEMORY INTERFACE

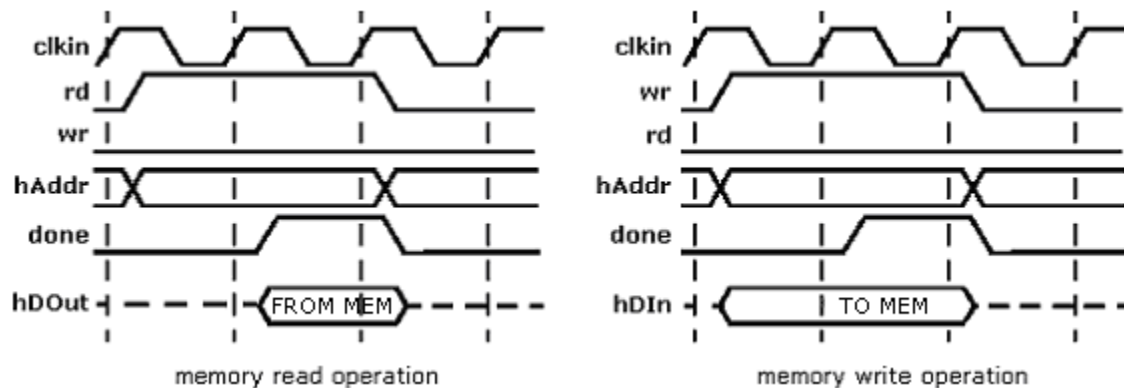
### INTERFACE:

The basic black box interface to memory is illustrated below. Note that the SDRAM controller (sdramcntl) is a black box component that cannot be changed, and its VHDL code will be provided to the student design teams. Therefore, only the interface to this subsystem is important.



### TIMING:

Below, the clock diagrams for both the memory read and memory write operations illustrate the basic behavior of the sdramcntl. The hAddr, hDIn, and hDOut lines are 16-bit busses. The remaining lines are single-bit standard logic.



---

# ADD

---

## ARITHMETIC INTEGER ADDITION

---

**OPCODE:**  
0000

**OPERATION:**  
 $\text{REG}[\text{DESTINATION}] = \text{REG}[\text{SOURCE ONE}] + \text{REG}[\text{SOURCE TWO}]$

**DESCRIPTION:**  
The register referenced by source one is added to the register referenced by source two, and the resulting value is placed in the register referenced by the destination. All 16 registers are addressable by both sources and the destination.

**INSTRUCTION WORD FORMAT:**

LOCALITY	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
GROUPING	OPCODE				DESTINATION				SOURCE ONE				SOURCE TWO			
VALUES	0	0	0	0	C	C	C	C	A	A	A	A	B	B	B	B

**ASSEMBLY CODE EXAMPLE:**  
ADD R7 R5 R6 # R7 = R5+R6

**BINARY EXAMPLE:**

INSTRUCTION	0000	0111	0101	0110
OPERATION	ADD	R7	R5	R6

**OPERAND REFERENCE:**

DESTINATION	
CCCC	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

SOURCE ONE	
AAAA	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

SOURCE TWO	
BBBB	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

---

# SUB

---

## ARITHMETIC INTEGER SUBTRACTION

---

**OPCODE:**

0001

**OPERATION:**

REG[DESTINATION] = REG[SOURCE ONE] - REG[SOURCE TWO]

**DESCRIPTION:**

The register referenced by source two is subtracted from the register referenced by source one, and the resulting value is placed in the register referenced by the destination. All 16 registers are addressable by both sources and the destination.

**INSTRUCTION WORD FORMAT:**

LOCALITY	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
GROUPING	OPCODE				DESTINATION				SOURCE ONE				SOURCE TWO			
VALUES	0	0	0	1	C	C	C	C	A	A	A	A	B	B	B	B

**ASSEMBLY CODE EXAMPLE:**

SUB R7 R5 R6 # R7 = R5-R6

**BINARY EXAMPLE:**

INSTRUCTION	0001	0111	0101	0110
OPERATION	SUB	R7	R5	R6

**OPERAND REFERENCE:**

DESTINATION	
CCCC	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

SOURCE ONE	
AAAA	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

SOURCE TWO	
BBBB	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3



---

# AND

---

## LOGICAL BITWISE AND

---

**OPCODE:**

0010

**OPERATION:**

REG[DESTINATION] = REG[SOURCE ONE] &amp; REG[SOURCE TWO]

**DESCRIPTION:**

The register referenced by source one is ANDed with the register referenced by source two, and the resulting value is placed in the register referenced by the destination. All 16 registers are addressable by both sources and the destination.

**INSTRUCTION WORD FORMAT:**

LOCALITY	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
GROUPING	OPCODE				DESTINATION				SOURCE ONE				SOURCE TWO			
VALUES	0	0	1	0	C	C	C	C	A	A	A	A	B	B	B	B

**ASSEMBLY CODE EXAMPLE:**

AND R7 R5 R6 # R7 = R5 AND R6

**BINARY EXAMPLE:**

INSTRUCTION	0010	0111	0101	0110
OPERATION	AND	R7	R5	R6

**OPERAND REFERENCE:**

DESTINATION	
CCCC	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

SOURCE ONE	
AAAA	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

SOURCE TWO	
BBBB	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

---

# OR

---

## LOGICAL BITWISE INCLUSIVE OR

---

**OPCODE:**

0011

**OPERATION:**

REG[DESTINATION] = REG[SOURCE ONE] | REG[SOURCE TWO]

**DESCRIPTION:**

The register referenced by source one is inclusively ORed with the register referenced by source two, and the resulting value is placed in the register referenced by the destination. All 16 registers are addressable by both sources and the destination.

**INSTRUCTION WORD FORMAT:**

LOCALITY	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
GROUPING	OPCODE				DESTINATION				SOURCE ONE				SOURCE TWO			
VALUES	0	0	1	1	C	C	C	C	A	A	A	A	B	B	B	B

**ASSEMBLY CODE EXAMPLE:**

OR R7 R5 R6 # R7 = R5 OR R6

**BINARY EXAMPLE:**

INSTRUCTION	0011	0111	0101	0110
OPERATION	OR	R7	R5	R6

**OPERAND REFERENCE:**

DESTINATION	
CCCC	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

SOURCE ONE	
AAAA	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

SOURCE TWO	
BBBB	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

# NAND

## LOGICAL BITWISE NAND

### OPCODE:

0100

### OPERATION:

REG[DESTINATION] = NOT(REG[SOURCE ONE] & REG[SOURCE TWO])

### DESCRIPTION:

The register referenced by source one is NANDed with the register referenced by source two, and the resulting value is placed in the register referenced by the destination. All 16 registers are addressable by both sources and the destination.

### INSTRUCTION WORD FORMAT:

LOCALITY	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
GROUPING	OPCODE				DESTINATION				SOURCE ONE				SOURCE TWO			
VALUES	0	1	0	0	C	C	C	C	A	A	A	A	B	B	B	B

### ASSEMBLY CODE EXAMPLE:

NAND R7 R5 R6 # R7 = R5 NAND R6

### BINARY EXAMPLE:

INSTRUCTION	0100	0111	0101	0110
OPERATION	NAND	R7	R5	R6

### OPERAND REFERENCE:

DESTINATION	
CCCC	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

SOURCE ONE	
AAAA	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

SOURCE TWO	
BBBB	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

# SLA

## ARITHMETIC SHIFT LEFT

**OPCODE:**

0101

**OPERATION:**

REG[DESTINATION] = REG[SOURCE] SLA VAL[OFFSET]

**DESCRIPTION:**

The register referenced by the source is arithmetically shifted left by the value defined in the offset, and the resulting value is placed in the register referenced by the destination. All 16 registers are addressable by the source and the destination.

**INSTRUCTION WORD FORMAT:**

LOCALITY	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
GROUPING	OPCODE				DESTINATION				SOURCE				OFFSET			
VALUES	0	1	0	1	C	C	C	C	A	A	A	A	B	B	B	B

**ASSEMBLY CODE EXAMPLE:**

SLA R7 R5 6 # R7 = R5 SLA 6

**BINARY EXAMPLE:**

INSTRUCTION	0101	0111	0101	0110
OPERATION	SLA	R7	R5	6 Bits

**OPERAND REFERENCE:**

DESTINATION	
CCCC	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

SOURCE	
AAAA	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

OFFSET	
BBBB	SHIFTED
0000	0 Bits
0001	1 Bits
0010	2 Bits
0011	3 Bits
0100	4 Bits
0101	5 Bits
0110	6 Bits
0111	7 Bits
1000	8 Bits
1001	9 Bits
1010	10 Bits
1011	11 Bits
1100	12 Bits
1101	13 Bits
1110	14 Bits
1111	15 Bits

# SRA

## ARITHMETIC SHIFT RIGHT

**OPCODE:**

0110

**OPERATION:**

REG[DESTINATION] = REG[SOURCE] SRA VAL[OFFSET]

**DESCRIPTION:**

The register referenced by the source is arithmetically shifted right by the value defined in the offset, and the resulting value is placed in the register referenced by the destination. All 16 registers are addressable by the source and the destination.

**INSTRUCTION WORD FORMAT:**

LOCALITY	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
GROUPING	OPCODE				DESTINATION				SOURCE				OFFSET			
VALUES	0	1	1	0	C	C	C	C	A	A	A	A	B	B	B	B

**ASSEMBLY CODE EXAMPLE:**

SRA R7 R5 6 # R7 = R5 SRA 6

**BINARY EXAMPLE:**

INSTRUCTION	0110	0111	0101	0110
OPERATION	SRA	R7	R5	6 Bits

**OPERAND REFERENCE:**

DESTINATION	
CCCC	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

SOURCE	
AAAA	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

OFFSET	
BBBB	SHIFTED
0000	0 Bits
0001	1 Bits
0010	2 Bits
0011	3 Bits
0100	4 Bits
0101	5 Bits
0110	6 Bits
0111	7 Bits
1000	8 Bits
1001	9 Bits
1010	10 Bits
1011	11 Bits
1100	12 Bits
1101	13 Bits
1110	14 Bits
1111	15 Bits

# LD

## LOAD FROM MEMORY

**OPCODE:**

0111

**OPERATION:**

REG[DESTINATION] ← MEM[SOURCE]

**DESCRIPTION:**

The memory location referenced by the source register is copied to the register referenced by the destination operand. The MODE must be set to 0110.

**INSTRUCTION WORD FORMAT:**

LOCALITY	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
GROUPING	OPCODE				DESTINATION				SOURCE				MODE			
VALUES	0	1	1	1	C	C	C	C	A	A	A	A	B	B	B	B

**ASSEMBLY CODE EXAMPLE:**

LD R7 R5 # R7 = MEM[R5]

**BINARY EXAMPLE:**

INSTRUCTION	0111	0111	0101	0110
OPERATION	LD/ST	R7	MEM[R5]	LD

**OPERAND REFERENCE:**

DESTINATION	
CCCC	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

SOURCE	
AAAA	REGISTER
0000	MEM[R0]
0001	MEM[R1]
0010	MEM[R2]
0011	MEM[R3]
0100	MEM[R4]
0101	MEM[R5]
0110	MEM[R6]
0111	MEM[R7]
1000	MEM[R8]
1001	MEM[R9]
1010	MEM[R10]
1011	MEM[R11]
1100	U0
1101	U1
1110	U2
1111	U3

MODE	
BBBB	FUNCTION
0110	LOAD
1001	STORE
others	NOP

# ST

## STORE TO MEMORY

**OPCODE:**

0111

**OPERATION:**

REG[SOURCE] → MEM[DESTINATION]

**DESCRIPTION:**

The register value referenced by the source operand is copied to the memory location referenced by the destination register. The MODE must be set to 1001.

**INSTRUCTION WORD FORMAT:**

LOCALITY	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
GROUPING	OPCODE				SOURCE				DESTINATION				MODE			
VALUES	0	1	1	1	C	C	C	C	A	A	A	A	B	B	B	B

**ASSEMBLY CODE EXAMPLE:**

ST R7 R5 # MEM[R5] = R7

**BINARY EXAMPLE:**

INSTRUCTION	0111	0111	0101	1001
OPERATION	LD/ST	R7	MEM[R5]	ST

**OPERAND REFERENCE:**

SOURCE	
CCCC	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

DESTINATION	
AAAA	REGISTER
0000	MEM[R0]
0001	MEM[R1]
0010	MEM[R2]
0011	MEM[R3]
0100	MEM[R4]
0101	MEM[R5]
0110	MEM[R6]
0111	MEM[R7]
1000	MEM[R8]
1001	MEM[R9]
1010	MEM[R10]
1011	MEM[R11]
1100	U0
1101	U1
1110	U2
1111	U3

MODE	
BBBB	FUNCTION
0110	LD
1001	ST
others	NOP

---

# LIL

---

## IMMEDIATE LOW-BYTE LOAD

---

**OPCODE:**

1000

**OPERATION:**

REG[DESTINATION][LOW] = VAL[IMMEDIATE]

**DESCRIPTION:**

The register referenced by the destination receives the value defined by the immediate portion of the instruction into its low byte. The contents of the destination register's high byte remain unchanged by this instruction.

**INSTRUCTION WORD FORMAT:**

LOCALITY	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
GROUPING	OPCODE				DESTINATION				IMMEDIATE							
VALUES	1	0	0	0	C	C	C	C	A	A	A	A	A	A	A	A

**ASSEMBLY CODE EXAMPLE:**

LIL R7 01010110 # R7[LOW] = 01010110

**BINARY EXAMPLE:**

INSTRUCTION	1000	0111	01010110
OPERATION	LIL	R7	0x56

**OPERAND REFERENCE:**

DESTINATION	
CCCC	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3



---

# LIH

---

## IMMEDIATE HIGH-BYTE LOAD

---

**OPCODE:**

1001

**OPERATION:**

REG[DESTINATION][HIGH] = VAL[IMMEDIATE]

**DESCRIPTION:**

The register referenced by the destination receives the value defined by the immediate portion of the instruction into its high byte. The contents of the destination register's low byte remain unchanged by this instruction.

**INSTRUCTION WORD FORMAT:**

LOCALITY	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
GROUPING	OPCODE				DESTINATION				IMMEDIATE							
VALUES	1	0	0	1	C	C	C	C	A	A	A	A	A	A	A	A

**ASSEMBLY CODE EXAMPLE:**

LIH R7 01010110 # R7[HIGH] = 01010110

**BINARY EXAMPLE:**

INSTRUCTION	1001	0111	01010110
OPERATION	LIH	R7	0x56

**OPERAND REFERENCE:**

DESTINATION	
CCCC	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

# CMP

## LOGICAL REGISTER DATA COMPARISON

**OPCODE:**

1010

**OPERATION:**

if (REG[SOURCE ONE] - REG[SOURCE TWO]) &lt; 0 (CRN = 1) else (CRN = 0)

if (REG[SOURCE ONE] - REG[SOURCE TWO]) == 0 (CRZ = 1) else (CRZ = 0)

if (REG[SOURCE ONE] - REG[SOURCE TWO]) &gt; 0 (CRP = 1) else (CRP = 0)

**DESCRIPTION:**

The register referenced by source two is compared to the register referenced by source one. All of the condition code registers are set for each comparison.

**INSTRUCTION WORD FORMAT:**

LOCALITY	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
GROUPING	OPCODE				UNUSED				SOURCE ONE				SOURCE TWO			
VALUES	1	0	1	0	X	X	X	X	A	A	A	A	B	B	B	B

**ASSEMBLY CODE EXAMPLE:**

CMP R5 R6 # R5 &gt;, =, or &lt; R6?

**BINARY EXAMPLE:**

INSTRUCTION	1010	XXXX	0101	0110
OPERATION	CMP	N/A	R5	R6

**OPERAND REFERENCE:**

SOURCE ONE	
AAAA	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

SOURCE TWO	
BBBB	REGISTER
0000	R0
0001	R1
0010	R2
0011	R3
0100	R4
0101	R5
0110	R6
0111	R7
1000	R8
1001	R9
1010	R10
1011	R11
1100	U0
1101	U1
1110	U2
1111	U3

# BRN / BRZ / BRP / JMP

## CONDITIONAL BRANCH OR JUMP

### OPCODE:

1011

### OPERATION:

```

if (MODE==1001)
    PC = (PC+1+VAL[OFFSET])
else if ([ (MODE==0110) AND (CRN==1) ] OR
        [ (MODE==0111) AND (CRZ==1) ] OR
        [ (MODE==1000) AND (CRP==1) ])
    PC = (PC+1+VAL[OFFSET])
else
    PC = (PC+1)

```

### DESCRIPTION:

If the MODE is JMP, the PC is set to the part of the program indicated in the assembly code by LABEL. (Note: the offset to reach LABEL is calculated by the assembler based on the incremented PC.) If the MODE is BRN, BRZ, or BRP, then the branch is conditional on CRN, CRZ, or CRP, respectively.

### INSTRUCTION WORD FORMAT:

LOCALITY	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
GROUPING	OPCODE				MODE				OFFSET							
VALUES	1	0	1	1	C	C	C	C	A	A	A	A	A	A	A	A

### ASSEMBLY CODE EXAMPLE:

BRN LABEL                      BRZ LABEL                      BRP LABEL                      JMP LABEL

### BINARY EXAMPLE:

INSTRUCTION	1011	0110	01010110
OPERATION	BR/JMP	BRN	0x56

### OPERAND REFERENCE:

MODE	
CCCC	TYPE
0110	BRN
0111	BRZ
1000	BRP
1001	JMP
others	NOP

# Appendix

## INSTRUCTION SET TARGET HARDWARE

### FPGA PROTOTYPING PLATFORM:

The system implementation platform is the fully integrated XESS XSA-100 FPGA prototyping board. This platform is equipped with a state of the art FPGA from the Xilinx Spartan II series (XC2S100). To complement this FPGA are a variety of digitally integrated systems, including a 16 MB SDRAM, a 256 KB Flash, a 100 MHz programmable (divisible by integers) oscillator, an integrated seven-segment display, dip switches, and a parallel port interface. These systems and interfaces will be explored in more detail throughout the semester.

The diagram below, courtesy of the XESS Corporation, illustrates some of the components and interfaces that will be utilized by students implementing this core. The important parts of this diagram are the Spartan II FPGA, the SDRAM, and the seven-segment display. These systems will be directly or indirectly manipulated via student synthesized VHDL code. More information pertaining to the target hardware may be found at: <http://www.xess.com/prod026.php3>.

