

Principles of Wireless Sensor Networks

<https://kth.instructure.com/courses/2912>

Lecture 3 **Physical Layer**

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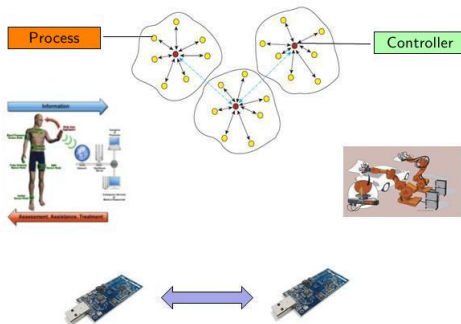
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Where we are



- How messages are successfully transmitted and received over the wireless channel?
- Aim: modeling the probability to successfully receive messages as function of the radio power, modulations, coding, and channel attenuations normally experienced in WSNs

Today's learning goals

- How bits of messages are transmitted over a channel?
- What is the probability to successfully receive messages over AWGN channels?
- What is the probability to successfully receive messages over fading channels?
- What is the Gilbert Eliot model for fading + AWGN channel + coding + modulation?

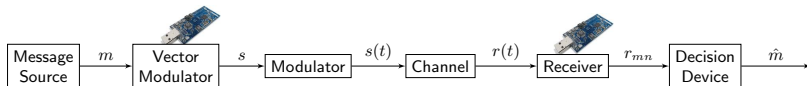
Outline

- Basics of modulation theory
- Probability of error over AWGN channels
- Probability of error over fading + AWGN channels
- Gilbert-Elliot model

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Digital modulations



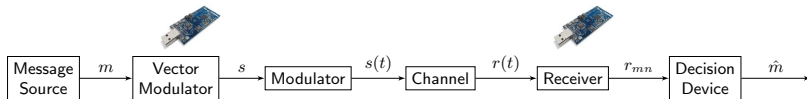
$$s(t) = \sum_{k=0}^{\infty} a_k(t) g(t - kT_s) \quad \text{modulated signal transmitted by the node's antenna}$$

$$G(f) \triangleq \int_0^{T_s} a_0 g(t) e^{-2\pi f t} dt \quad \text{spectrum of the signal over a symbol}$$

$$\Phi_s(f) \triangleq \frac{|G(f)|^2}{T_s} \quad \text{power spectral density of the signal}$$

$$T_s \quad \text{symbol duration}$$

Example: Binary Phase Shift Keying (BPSK) modulation



$$s(t) = \sum_{k=0}^{\infty} a_k(t) g(t - kT_s)$$

$$a_k(t) = \begin{cases} \cos(2\pi f_c t) & \text{if bit 1 at symbol time } k, \\ \cos(2\pi f_c t + \pi) & \text{if bit 0 at symbol time } k. \end{cases} \quad \text{Modulation of the bits}$$

$$g(t) = \sqrt{\frac{E}{T_s}}, \quad 0 \leq t \leq T_s \quad \quad P_t = \frac{E}{T_s} \quad \text{transmit power}$$

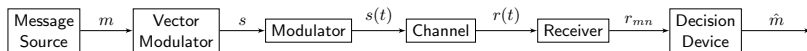
Probability of successful bit reception

Now, we would like to compute the probability that a bit is received successfully when it is transmitted by a modulation over an AWGN channel

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BPSK probability of error in AWGN channels

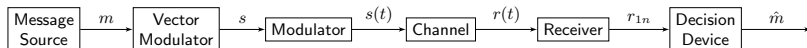


- Assume that the transmitted signal is received corrupted by an Additive White Gaussian Noise (AWGN)
- Assume that there is no fading, namely $A = 1$

$$r(t) = s(t) + n_0(t)$$

$$n_0(t) \in N\left(0, \sigma^2 = \frac{N_0}{2T_s}\right)$$

BPSK detection in AWGN wireless channels



- After a matched filter, the demodulator in the receiver produces a signal

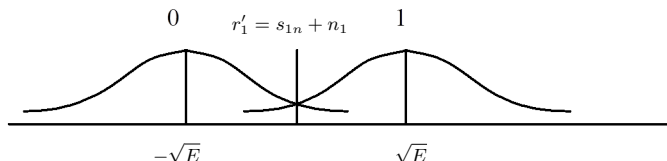
$$r'_1 = s_{1n} + n_1$$

$$s_{1n} = \begin{cases} \sqrt{E} & \text{if bit 1 was transmitted} \\ -\sqrt{E} & \text{if bit 0 was transmitted.} \end{cases}$$

$$n_1 \in \mathcal{N}\left(0, \sigma^2 = \frac{N_0}{2}\right)$$

- If $r_1 \geq 0$ the detector decides for bit 1
- If $r_1 < 0$ the detector decides for bit 0
- Given the AWGN, what is the error in this detection?

BPSK probability of error

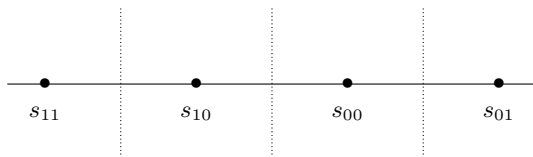


$$P_{e,0|1} \triangleq \int_{-\infty}^0 \frac{1}{\sqrt{2\pi\sigma^2}} e^{-\frac{(x-\sqrt{E})^2}{2\sigma^2}} dx = Q\left(\sqrt{\frac{2E}{N_0}}\right)$$

$$P_{e,1|0} \triangleq \int_0^{\infty} \frac{1}{\sqrt{2\pi\sigma^2}} e^{-\frac{(x+\sqrt{E})^2}{2\sigma^2}} dx = Q\left(\sqrt{\frac{2E}{N_0}}\right)$$

$$P_e \triangleq \frac{1}{2}P_{e,0|1} + \frac{1}{2}P_{e,1|0} = Q\left(\sqrt{\frac{2E}{N_0}}\right) \quad \text{SNR} \triangleq \frac{2E}{N_0}$$

Amplitude modulation, AM



- BPSK is a simple example of binary Amplitude Modulation
- By adding more amplitude values, more general modulations are possible
- Modulations are characterized by constellation points
- Every point (symbol) is associated to a signal with specific amplitude
- Every signal is then associated to a code-word of bits

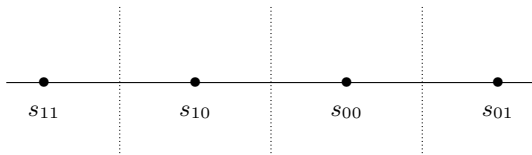
Probability of error for AM

- The distance between the constellation points determines the probability that a symbol is detected with error
- It is possible to show that

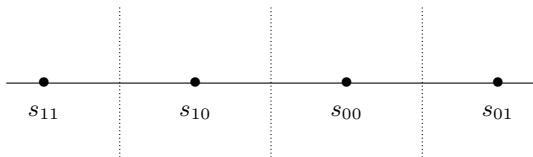
$$P_e \simeq N_{d_{\min}} Q\left(\frac{d_{\min}}{\sqrt{2N_0}}\right)$$

d_{\min} minimum distance among points

$N_{d_{\min}}$ average number of neighbors at minimum distance



Exercise: 4-AM (or 4-PAM)



$$s_{00} = \sqrt{\frac{E}{5}}g(t) \quad s_{01} = 3\sqrt{\frac{E}{5}}g(t)$$

$$s_{11} = -3\sqrt{\frac{E}{5}}g(t) \quad s_{10} = -\sqrt{\frac{E}{5}}g(t)$$

$$g(t) = \sqrt{\frac{1}{T_s}} \quad 0 \leq t \leq T_s$$

What is the probability that a bit is erroneously received?

Solution

$$d_{\min}^2 = 4\frac{E}{5} \quad N_{\min} = 1.5$$

- The probability of symbol error is

$$P_e \simeq 1.5Q\left(\sqrt{\frac{2E}{5N_0}}\right) = 1.5Q\left(\sqrt{\frac{4E_b}{5N_0}}\right)$$

- Since per every symbol there are two bits, the probability of a bit in error is

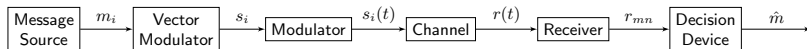
$$P_b \simeq 0.75Q\left(\sqrt{\frac{4E_b}{5N_0}}\right)$$

$$E_b = \frac{E}{\log_2 M} = \frac{E}{2} \quad \text{Energy per bit}$$

4-Quadrature Amplitude Modulation

- Let's consider a more general amplitude modulation that is used in
 - ▶ TskyMotes
 - ▶ IEEE 802.15.4 standard (which we study next lecture)
- 4-Quadrature Amplitude Modulation - 4QAM

4-QAM

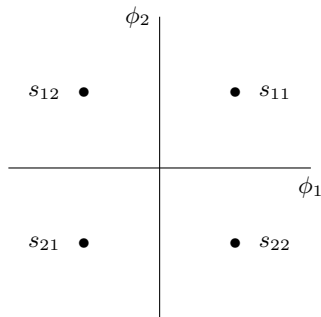


$$\begin{aligned}
 s_n(t) &= \frac{2E}{T_s} \cos \left(2\pi f_c t + \frac{(2n-1)\pi}{4} \right) \quad 0 \leq t \leq T_s \quad n = 1, \dots, 4 \\
 &= \frac{2E}{T_s} \cos \left(\frac{(2n-1)\pi}{4} \right) \cos(2\pi f_c t) - \frac{2E}{T_s} \sin \left(\frac{(2n-1)\pi}{4} \right) \sin(2\pi f_c t) \quad 0 \leq t \leq T_s
 \end{aligned}$$

Signal space by two basis functions

$$\phi_1(t) = \frac{2}{T_s} \cos(2\pi f_c t) \quad 0 \leq t \leq T_s$$

$$\phi_2(t) = \frac{2}{T_s} \sin(2\pi f_c t) \quad 0 \leq t \leq T_s$$



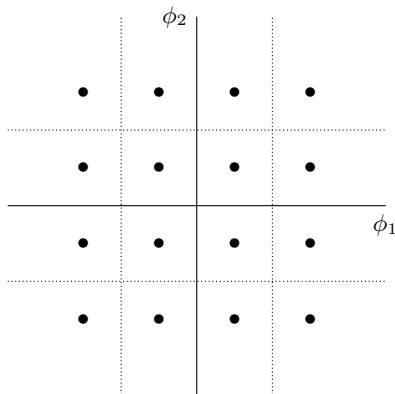
Probability of error in 4-QAM

Minimum distance between two symbols $d_{\min} = \sqrt{2E}$

2 neighbors at minimum distance

$$P_e \simeq 2Q\left(\sqrt{\frac{E}{5N_0}}\right) = 2Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$$

A more complex modulation: 16-QAM



More points can be added, this increases the transmit bit rate, but increases also the probability of error

Comparison of probabilities of error

$$P_{e,\text{BPSK}} = Q\left(\sqrt{\frac{2E}{N_0}}\right) \qquad \text{SNR} = \frac{2E}{N_0}$$

$$P_{e,4\text{-PAM}} \simeq 1.5Q\left(\sqrt{\frac{2E}{5N_0}}\right)$$

$$P_{e,4\text{-QAM}} \simeq 2Q\left(\sqrt{\frac{E}{N_0}}\right)$$

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Communication over wireless channel

- So far, we have seen only channels where the transmitted signal is received corrupted by AWGN, but in real channel

$$r(t) = \sqrt{A}s(t) + n_0(t)$$

- The power of the transmitted signals is attenuated by the wireless channel
- How the bit probability of error is affected by fading channels?

Probability of error over fading channels

- Consider a AWGN + Rayleigh channel with path loss, Rayleigh fast fading, and fixed shadow fading

$$P_r = P_t G_t(\theta_t, \psi_t) G_r(\theta_r, \psi_r) \underbrace{\frac{\lambda^2}{(4\pi r)^2} \overline{\text{PL}} \cdot y \cdot z}_{\triangleq P_t C \cdot z}$$

Probability of error over fading channels

- The receiver sees the transmit power $P_t = E/T_s$ received with an attenuation $C \cdot z$
- We can reuse the probability for AWGN channels AS if there WERE no fading and the transmit power WERE

$$P_t = \frac{E}{T_s} C z$$

- Thus, the probability of error with fading has similar expression of simple AWGN channels, but with the SNR of a fading channel

$$\text{SNR} = \frac{2E}{N_0} C z$$

$$P_{e,\text{BPSK}} = Q\left(\sqrt{\frac{2E}{N_0}}\right) \longrightarrow P_{e,\text{BPSK}}(z) = Q\left(\sqrt{\frac{2ECz}{N_0}}\right)$$

AWGN no fading

AWGN + Rayleigh fading

Probability of error over fading channels

- The probability so derived over fading channel is instantaneous
 - ▶ Depends on the given realization of the fading channel z
- What is the average probability of error, where the average is taken over the distribution of the fading?
- Just take the expectation of $P_{e,\text{BPSK}}(z)$ over the distribution of z
- Thus (remember from Exercise session that the square of a Rayleigh random variable is an exponential random variable)

$$p(z) = \frac{1}{\gamma^*} e^{-\frac{z}{\gamma^*}} \quad \gamma^* \triangleq \mathbb{E}\{\text{SNR}\} = \frac{2E}{N_0} C \quad \mathbb{E}\{z\} = 1$$

$$\bar{P}_{e,\text{BPSK}} = \int_0^{\infty} P_{e,\text{BPSK}}(z) p(z) dz = \frac{1}{2} \left[1 - \sqrt{\frac{\gamma^*}{1 + \gamma^*}} \right] \simeq \frac{1}{4\gamma^*}$$

Error probability comparison over fading channels

- Simple AWGN channel:

$$P_{e,\text{BPSK}} = Q\left(\sqrt{\frac{2E}{N_0}}\right)$$

- ▶ Linear increase of SNR results in exponential decrease in the error probability

- AWGN + Rayleigh channel:

$$\bar{P}_{e,\text{BPSK}} \simeq \frac{1}{4\gamma^*} \quad \gamma^* \triangleq \mathbb{E}\{\text{SNR}\} = \frac{2E}{N_0}C$$

- ▶ Linear increase of the SNR gives only a linear decrease of error probability
- ▶ 40dB higher SNR than the simple AWGN channel to have same probability $P_e = 10^{-6}$

Packet error probability

- Bits are grouped in units denoted physical layer data frames, “physical layer messages”
- What is the probability that the message is erroneously received?
- Suppose that the message is formed by L bits and BPSK is used, then the probability that the packet is in error is bounded by

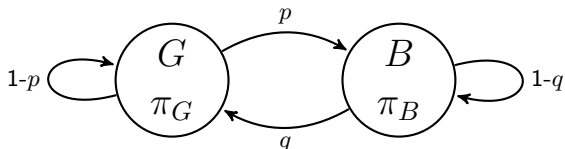
$$P_m \leq 1 - (1 - P_{e,\text{BPSK}})^L$$

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Gilbert-Elliot model

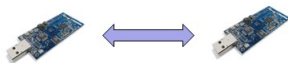
- It is a simple way to describe the behavior of the wireless channel in two states: Bad and Good



$$\pi_B + \pi_G = 1$$
$$\pi_G = (1 - p) \pi_G + q \pi_B$$

- π_B probability of bad state
- π_G probability of good state
- p probability to go from the good state to the bad
- q probability to go from the bad state to the good
- We will see in the exercise session how to tie these probabilities to the message size and the wireless channel

Conclusions



- We studied how bits are modulated and transmitted over the wireless channel
- The probability of successful reception of messages was characterized for AWGN and for AWGN + fading channels

Next lecture

- But how a node has the right to transmit a message?
- Medium access control
 - ▶ When a node gets the right to transmit?
 - ▶ What is the mechanism to get such a right?